Chapter 3, Case Studies:

Linear-time sorting algorithms (last week Wednesday and Friday)

◆□▶ ◆□▶ ◆注▶ ◆注▶ 注 のへで

- Disjoint sets and array forests (Today)
- Priority queues and heaps (Wednesday and Friday)
- N-sets and bit vectors (Thursday lab)
- (Begin Graph unit in lab next week)

Today:

- Quiz solutions
- Problem statement
- Disjoint set ADT details
- The array forest abstraction and data structure
- Find and union strategies, with optimizations

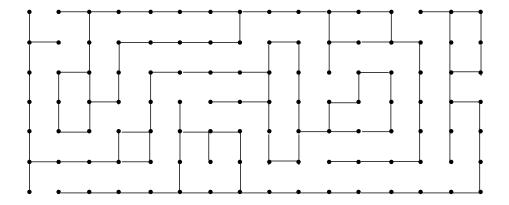
```
static Node arrayToList1(int[] array) {
    Node toReturn = new Node(array[0], null);
    for (int i = 1; i < array.length; i++) {</pre>
        Node current = toReturn;
        while (current.next() != null)
            current = current.next():
        current.setNext(new Node(array[i], null));
    3
    return toReturn:
}
Node arrayToList2(int[] array) {
    Node toReturn = null:
    for (int i = array.length - 1; i >= 0; i--)
        toReturn = new Node(array[i], toReturn);
    return toReturn:
}
static int[] listToArray(Node head) {
    int size = 0;
    for (Node current = head: current != null: current = current.next())
        size++:
    int[] toReturn = new int[size];
    int i = 0:
    for (Node current = head; current != null; current = current.next())
        toReturn[i++] = current.datum();
    return toReturn:
}
```

```
◆□▶ ◆□▶ ◆三▶ ◆三▶ 三三 のへで
```

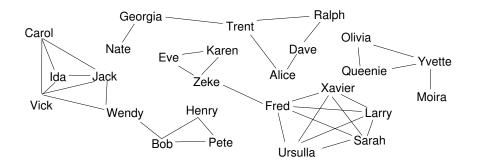
Problem statement:

Suppose we have a collection of items connected by an unknown equivalence relation. Efficiently find the equivalence classes in this collection as information about the relation is discovered.

◆□▶ ◆□▶ ◆臣▶ ◆臣▶ 臣 -



▲□▶ ▲圖▶ ▲目▶ ▲目▶ 目 のへで



・ロト ・個ト ・モト ・モト

2

a = c e = a + b d = b g = 1 f = d + ch = e * g

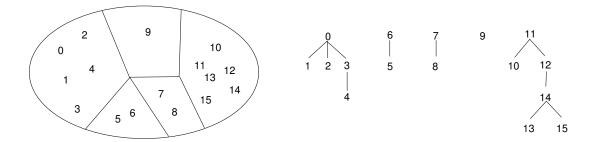
◆□▶ ◆□▶ ◆臣▶ ◆臣▶ 臣 の�?

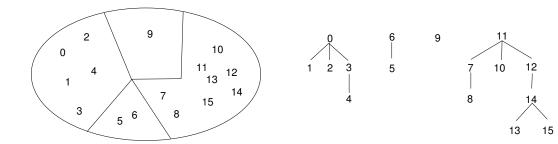
The disjoint set ADT:

Main operations: union two sets, find a set for a given element, and test if two elements are in the same set.

▲ロト ▲圖ト ▲画ト ▲画ト 三回 - のへで

- The universe is closed.
- We assume all elements can be indexed, [0, N).
- A set in the partition is identified by a leader.



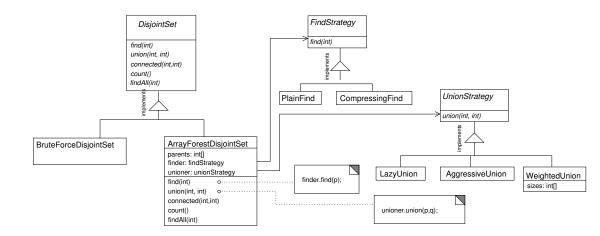


Invariant (Class ArrayForestDisjointSet)

For all $i \in [0, n)$,

- (a) leader(i) = leader(parents[i]), that is, parents[i] points to another element in the same set as i.
- (b) leader(i) = parents[leader(i)], that is, leaders all point to themselves.
- (c) Following a finite number links implied by parents will converge, that is, there is no circularity in the tree.

◆□▶ ◆□▶ ◆臣▶ ◆臣▶ 臣 -



▲□▶ ▲□▶ ▲三≯ ▲三≯ ▲□▶

Union strategy	LazyUnion	AggressiveUnion	WeightedUnion	LazyUnion	WeightedUnion
Find strategy	PlainFind	PlainFind	PlainFind	CompressingFind	CompressingFind
Find heavy:	1.30E7	3.34E7	7.40E5	9.26E5	6.68E5
	(5.68E6)	(8.40E3)	(1.80E4)	(2.38E4)	(9.34E3)
Even mix:	9.89E7	4.41E7	1.20E6	1.56E6	9.80E5
	(1.22E7)	(9.93E3)	(1.97E4)	(2.12E4)	(9.96E3)
Union heavy:	1.62E8	4.39E7	1.40E6	1.71E6	1.04E6
	(1.26E7)	(9.99E3)	(2.01E4)	(1.59E4)	(1.00E4)

▲□▶ ▲□▶ ▲三▶ ▲三▶ 三三 - のへで

Coming up: (all end-of-day)

Do linear sorting project (Wed, Sept 23)

Due **Today**: Finish reading Section 3.2 (disjoint sets and array forests) Do Ex 2.(12 & 16) and 3.8 Take disjoint-sets quiz

▲□▶ ▲圖▶ ▲필▶ ▲필▶ 三里

Due **Thurs, Sept 26**: Read Section 3.4 Do Exercises 3.(26 & 27). Take N-sets quiz

Due Fri, Sept 27: Read Section 3.3 (heaps and priority queues) (no exercises) Take heap/pq quiz