Chapter 5, Binary search trees:

 Binary search trees; the balanced BST problem (spring-break eve; finished last week Monday)

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- AVL trees (last week Monday and Wednesday)
- Traditional red-black trees (last week Friday, finish Today)
- Left-leaning red-black trees (Today)
- "Wrap-up" BST (Wednesday)
- Begin dynamic programming (Friday)

Today:

- Finishing Traditional RB
- LLRB context and definition
- LLRB invariant and cases
- Performance comparison among AVL, TrRB, and LLRB

Why invariants?

- An invariant is a constraint we put on our code to help us guarantee something about it.
- ▶ The general invariant for BSTs guarantees the correctness of our find algorithm.
- The invariants for AVL trees and RB trees guarantee logarithmic-time operations. A stronger constraint is both a stronger constraint to *maintain* and a stronger constraint to *assume*.

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A left-leaning red-black tree is a binary tree (usually a BST) that is either empty or it is rooted at node T such that

- ► *T* is either red or black.
- ▶ Both of *T*'s children are roots of left-leaning red-black trees.
- ► *T*'s right child is black.
- ▶ If *T* is red, then its left child is black.
- The left-leaning red-black trees rooted at its children have equal blackheight; moreover, the blackheight of the tree rooted at T is one more that the blackheight of its children if T is black or equal to that of its children if T is red.

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Potential violations

Ignorant node

Inconsistent backheight

Red null

Double red

Right red

shouldn't happen

fix when they happen

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Invariant 28 (Postconditions of RealNode.put() with LLRBBalancer.) Let x be the root of a subtree on which put() is called and let y be the node returned, that is, the root of the resulting subtree.

- (a) The subtree rooted at y has a consistent black height.
- (b) The black height of subtree rooted at y is equal to the original black height of the subtree rooted at x.
- (c) The subtree rooted at y has no double-red violations except, possibly, both y and its left child is red, which can happen only if x is a left child.

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(d) The subtree rooted at y has no right-red violations.











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	After puts			After removals		
	Height	Leaf %	Total depth	Height	Leaf %	Total depth
Unbalanced	32	33.3%	134507	28	16.8%	61207
	31	33.2%	127865	26	17.0%	58171
	30	33.1%	129037	26	16.9%	58610
	28	33.5%	124463	26	17.3%	56086
	32	33.4%	136730	28	16.9%	62092
AVL	16	43.2%	100327	14	21.5%	46088
	15	42.9%	100395	14	21.1%	46028
	15	42.8%	100341	14	21.1%	46028
	15	42.8%	100282	14	21.3%	45973
	15	43.0%	100582	14	21.2%	46097
Traditional RB	16	42.8%	101948	16	21.5%	46729
	16	42.9%	101226	15	21.4%	46344
	16	43.1%	101525	15	21.5%	46462
	16	42.7%	101680	16	21.5%	46572
	16	42.9%	101292	15	21.4%	46338
Left-leaning RB	18	42.8%	102288	18	21.6%	46950
-	19	42.9%	102860	16	21.3%	46774
	18	43.1%	101949	17	21.5%	46691
	18	42.7%	102011	17	21.6%	46938
	19	42.9%	102552	16	21.4%	46764

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