

## Chapter 5, Dynamic Programming:

- ▶ Introduction and sample problems (previous week Wednesday)
- ▶ Principles of DP (previous week Friday)
- ▶ DP algorithms, solutions to sample problems (last week Monday)
- ▶ Introduce optimal BSTs / review for test 2 (last week Wednesday)
- ▶ **Test 2**, *not* covering DP (last week Friday)
- ▶ Retrospective on Test 2 (Monday)
- ▶ Finish up optimal BSTs (**Today**)
- ▶ [Begin hash tables (Friday)]

## Today:

- ▶ Review optimal BST definition
- ▶ The optimal-BST-building problem
- ▶ The dynamic programming solution

## Why this problem?

- ▶ It connects dynamic programming with the quest for a better map.
- ▶ Its hardness is in the right places (building the table—hard; reconstructing solution—trivial).
- ▶ It is a representative of a bigger concept: What if we had more information—how would that change the problem.

## Game plan:

- ▶ Understand the problem itself
- ▶ Understand the recursive characterization
- ▶ Understand the table-building algorithm

The **optimal binary search tree** problem:

- ▶ Assume we know all the keys  $k_0, k_1, \dots, k_{n-1}$  ahead of time.
- ▶ Assume further that we know the probabilities  $p_0, p_1, \dots, p_{n-1}$  of each key's lookup.
- ▶ Assume even further that we know the “miss probabilities”  $q_0, q_1, \dots, q_n$  where  $q_i$  is the probability that an *extraneous* key falling between  $k_{i-1}$  and  $k_i$  will be looked up.
- ▶ We want to build a tree to minimize the *expected cost* of a look up, which is the *total weighted depth* of the tree:

$$\sum_{i=0}^{n-1} p_i \text{ depth}(k_i) + \sum_{i=0}^n q_i \text{ depth}(m_i)$$

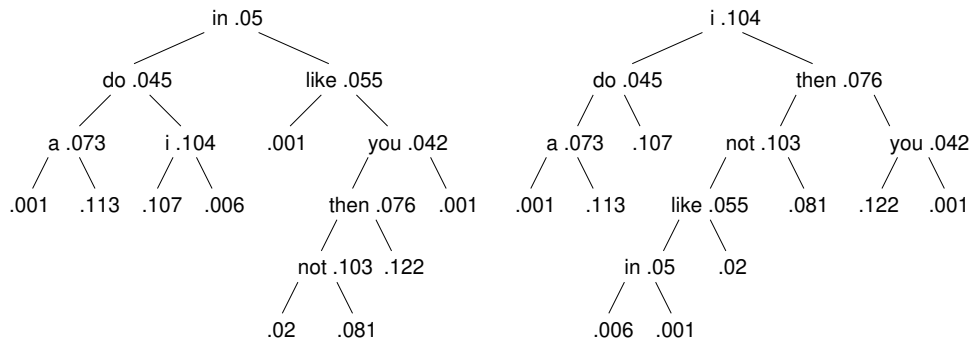
where  $\text{depth}(x)$  is the number of nodes to be inspected on the route from the root to node  $x$ ,  $k_i$  stands for the node containing key  $k_i$  [notational abuse], and  $m_i$  is the dummy node between keys  $k_{i-1}$  and  $k_i$ .

- ▶ Note that the rules of probability require  $\sum_{i=0}^{n-1} p_i + \sum_{i=0}^n q_i = 1$

i	84	eat	24	ham	10	fox	7	rain	4
not	83	will	21	there	9	on	7	see	4
them	61	sam	19	train	9	tree	6	try	4
a	59	with	19	anywhere	8	say	5	boat	3
like	44	am	16	house	8	so	5	that	3
in	40	could	14	mouse	8	be	4	are	2
do	36	here	11	or	8	goat	4	good	2
you	34	the	11	box	7	let	4	thank	2
would	26	eggs	10	car	7	may	4	they	2
and	24	green	10	dark	7	me	4	if	1

	Key or miss event	combined frequency
	{ }	0
	a	59
{	am and anywhere are be boat box car could dark }	92
	do	36
{	eat eggs fox goat good green ham here house }	86
	i	84
	{ if let }	5
	in	40
	{ }	0
	like	44
	{ may me mouse }	16
	not	83
{	on or rain same say see so thank that the }	65
	then	61
{	there they train tree try will with would }	99
	you	34
	{ }	0

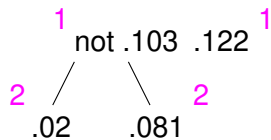
	0	1	2	3	4	5	6	7
$k_i$	a	do	i	in	like	not	then	you
$p_i$	.073	.045	.104	.05	.055	.103	.076	.042
$q_i$	.001	.113	.107	.006	.001	.02	.081	.122



$$1 \cdot .02 + 1 \cdot .081 \\ = .101$$

$$\begin{array}{r} 1 \\ .02 \\ 1 \\ .081 \end{array}$$

$$\begin{aligned} &2 \cdot .02 + 2 \cdot .081 \\ &+ 1 \cdot .103 + 1 \cdot .122 \\ &= .427 \end{aligned}$$



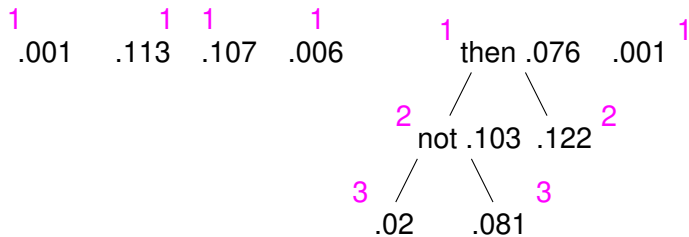


$$3 \cdot .02 + 3 \cdot .081$$

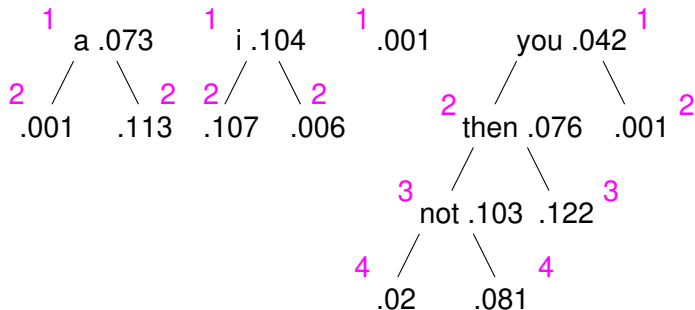
$$+ 2 \cdot .103 + 2 \cdot .122$$

$$+ 1 \cdot .001 + 1 \cdot .133 + 1 \cdot .107 + 1 \cdot .006 + 1 \cdot .076 + 1 \cdot .001$$

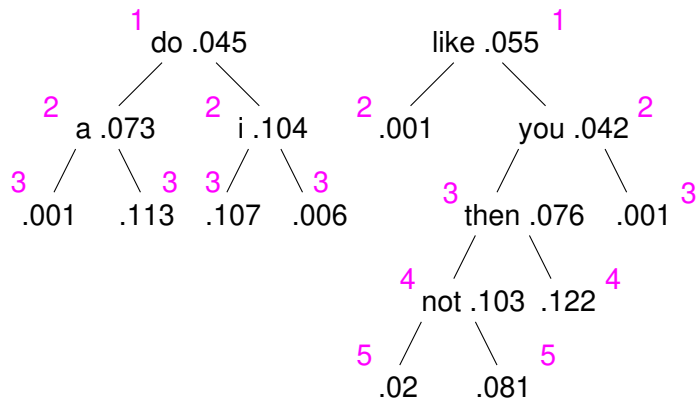
$$= 1.057$$



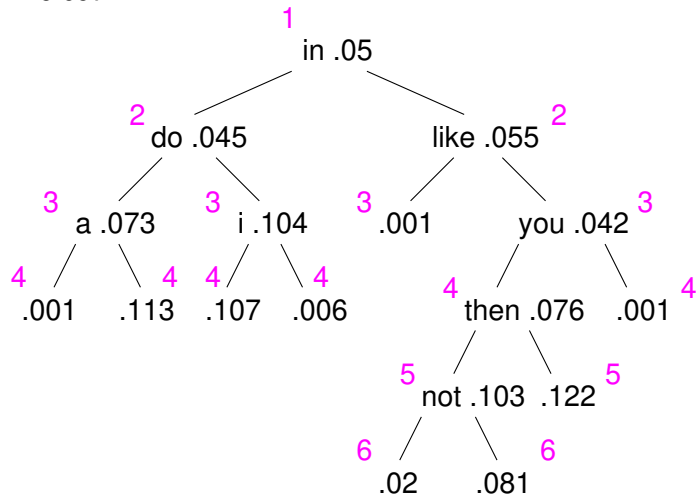
$$\begin{aligned}
&4 \cdot .02 + 4 \cdot .081 \\
&+ 3 \cdot .103 + 3 \cdot .122 \\
&+ 2 \cdot .001 + 2 \cdot .133 + 2 \cdot .107 + 2 \cdot .006 + 2 \cdot .076 + 2 \cdot .001 \\
&+ 1 \cdot .073 + 1 \cdot .104 + 1 \cdot .001 + 1 \cdot .042 \\
&= 1.907
\end{aligned}$$



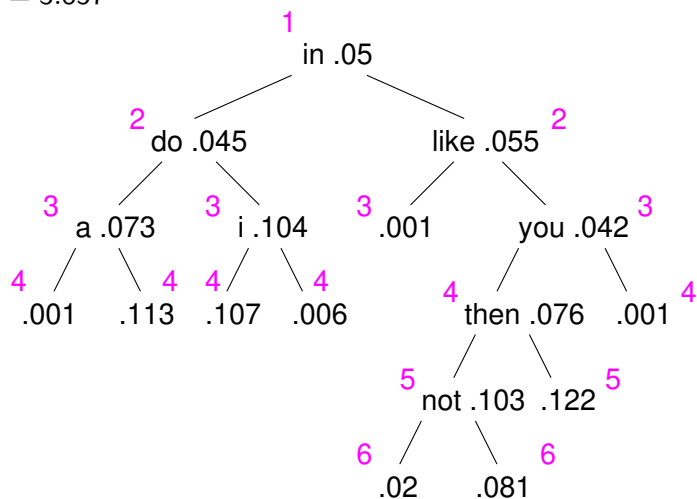
$5 \cdot .02 + 5 \cdot .081$   
 $+ 4 \cdot .103 + 4 \cdot .122$   
 $+ 3 \cdot .001 + 3 \cdot .133 + 3 \cdot .107 + 3 \cdot .006 + 3 \cdot .076 + 3 \cdot .001$   
 $+ 2 \cdot .073 + 2 \cdot .104 + 2 \cdot .001 + 2 \cdot .042$   
 $+ 1 \cdot .045 + 1 \cdot .055$   
 $= 2.857$



$6 \cdot .02 + 6 \cdot .081$   
 $+ 5 \cdot .103 + 5 \cdot .122$   
 $+ 4 \cdot .001 + 4 \cdot .133 + 4 \cdot .107 + 4 \cdot .006 + 4 \cdot .076 + 4 \cdot .001$   
 $+ 3 \cdot .073 + 3 \cdot .104 + 3 \cdot .001 + 3 \cdot .042$   
 $+ 2 \cdot .045 + 2 \cdot .055$   
 $+ 1 \cdot .05$   
 $= 3.857$

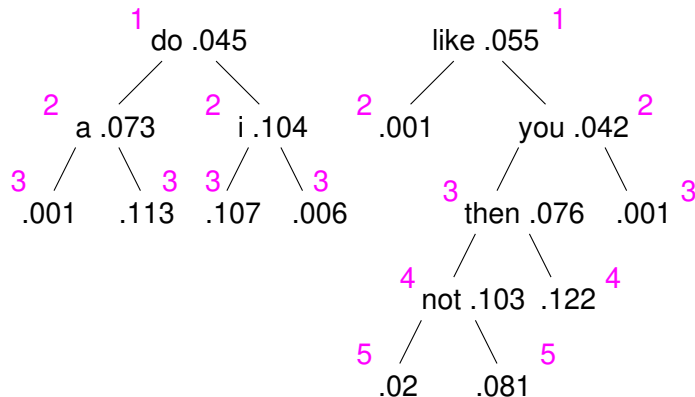


$$\begin{aligned}
 &4 \cdot .001 + 3 \cdot .073 + 4 \cdot .133 + 2 \cdot .045 + 4 \cdot .107 + 3 \cdot .104 + 4 \cdot .006 \\
 &+ 1 \cdot .05 \\
 &+ 3 \cdot .001 + 2 \cdot .055 + 6 \cdot .02 + 6 \cdot .081 + 4 \cdot .076 + 5 \cdot .122 + 3 \cdot .042 + 4 \cdot .001 \\
 &= 3.857
 \end{aligned}$$



$3 \cdot .001 + 2 \cdot .073 + 3 \cdot .133 + 1 \cdot .045 + 3 \cdot .107 + 2 \cdot .104 + 3 \cdot .006$   
 $+ 2 \cdot .001 + 1 \cdot .055 + 5 \cdot .02 + 5 \cdot .081 + 3 \cdot .076 + 4 \cdot .122 + 2 \cdot .042 + 3 \cdot .001$   
 $+ .001 + .073 + .133 + .045 + .107 + .104 + .006$   
 $+ .05$   
 $+ .001 + .055 + .02 + .081 + .076 + .122 + .042 + .001$   
 $= 3.857$

in .05



Total weighted depth for a given tree (expected lookup cost):

$$\underbrace{\sum_{i=0}^{n-1} p_i \text{depth}(k_i)}_{\text{keys}} + \underbrace{\sum_{i=0}^n q_i \text{depth}(m_i)}_{\text{misses}}$$

Let  $\text{depth}_{k_a}(k_i)$  be the depth of the node with  $k_i$  in the subtree rooted at node with  $k_a$ . For example, if  $k_r$  is the root of the entire tree and  $k_a$  is a child of the root, then

$$\text{depth}_{k_r}(k_i) = \text{depth}_{k_a}(k_i) + 1$$

Then we can rewrite the total weighted depth as

$$\underbrace{\sum_{i=0}^{r-1} p_i \text{depth}_{k_r}(k_i) + \sum_{i=0}^r q_i \text{depth}_{k_r}(m_i) + p_r}_{\text{left subtree total weighted depth (absolute)}} + \underbrace{\sum_{i=r+1}^{n-1} p_i \text{depth}_{k_r}(k_i) + \sum_{i=r+1}^n q_i \text{depth}_{k_r}(m_i)}_{\text{right subtree total weighted depth (absolute)}}$$

Again, let  $k_r$  be the root of the entire tree and  $k_a$  and  $k_b$  be the root's children. Then

$$\underbrace{\sum_{i=0}^{r-1} p_i(\text{depth}_{k_a}(k_i) + 1) + \sum_{i=0}^r q_i(\text{depth}_{k_a}(m_i) + 1) + p_r}_{\text{left subtree total weighted depth (absolute)}} + \underbrace{\sum_{i=r+1}^{n-1} p_i(\text{depth}_{k_b}(k_i) + 1) + \sum_{i=r+1}^n q_i(\text{depth}_{k_r}(m_i) + 1)}_{\text{right subtree total weighted depth (absolute)}}$$

Convert to "relative depth":

$$\underbrace{\sum_{i=0}^{n-1} p_i + \sum_{i=0}^n q_i}_{\text{total probability}} + \underbrace{\sum_{i=0}^{r-1} p_i \text{ depth}_{k_a}(k_i) + \sum_{i=0}^r q_i \text{ depth}_{k_a}(m_i)}_{\text{left subtree total weighted depth (relative)}} + \underbrace{\sum_{i=r+1}^{n-1} p_i \text{ depth}_{k_b}(k_i) + \sum_{i=r+1}^n q_i \text{ depth}_{k_r}(m_i)}_{\text{right subtree total weighted depth (relative)}}$$

Let  $TWD(k)$  be the total weighted depth of the tree rooted at  $k$  (relative to  $k$ ) and  $TP(k)$  be the total probability of the tree rooted at  $k$ . Then

$$TWD(k_r) = TP(k_r) + TWD(k_a) + TWD(k_b)$$



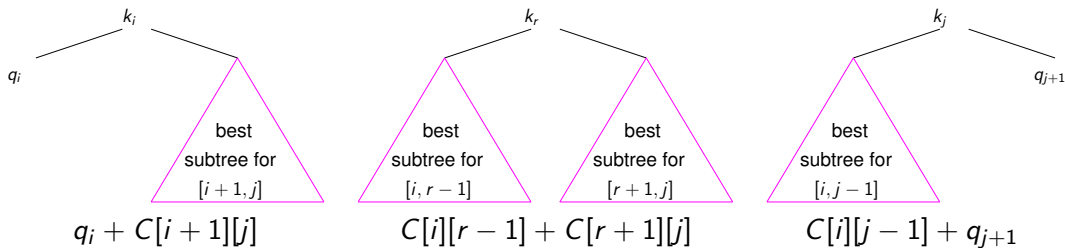
Let  $P[i][j]$  be the total probabilities of the keys and misses in the range  $[i, j]$ :

$$P[i][j] = \sum_{k=i}^j p_k + \sum_{k=i}^{j+1} q_k$$

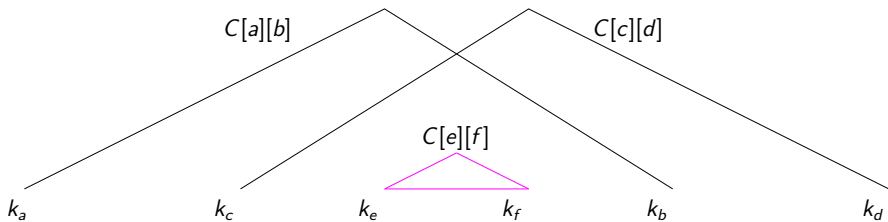
Let  $C[i][j]$  be the least total weighted depth of any BST composed from keys in the range  $[i, j]$ . The recursive characterization is

$$C[i][j] = \begin{cases} 2q_i + p_i + 2q_{i+1} & \text{if } i = j \\ P[i][j] + \min \left\{ \begin{array}{l} q_i + C[i+1][j] \\ C[i][r-1] + C[r+1][j] \text{ for } r \in (i, j) \\ C[i][j-1] + q_{j+1} \end{array} \right\} & \text{if } i < j \end{cases}$$

$$C[i][j] = \begin{cases} 2q_i + p_i + 2q_{i+1} & \text{if } i = j \\ P[i][j] + \min \left\{ \begin{array}{l} q_i + C[i+1][j] \\ C[i][r-1] + C[r+1][j] \text{ for } r \in (i, j) \\ C[i][j-1] + q_{j+1} \end{array} \right\} & \text{if } i < j \end{cases}$$



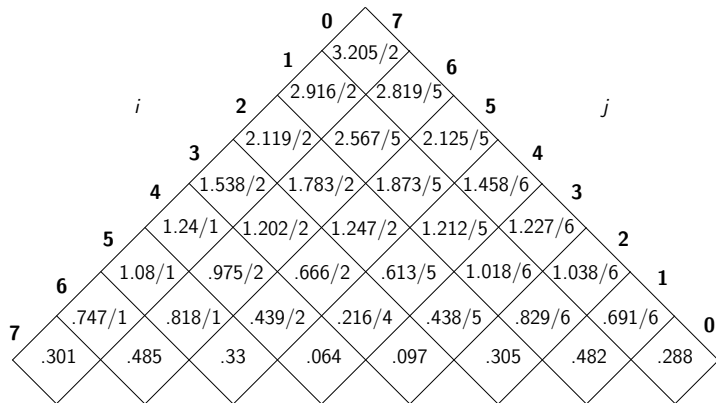
$$C[i][j] = \begin{cases} 2q_i + p_i + 2q_{i+1} & \text{if } i = j \\ P[i][j] + \min \left\{ \begin{array}{l} q_i + C[i+1][j] \\ C[i][r-1] + C[r+1][j] \text{ for } r \in (i, j) \\ C[i][j-1] + q_{j+1} \end{array} \right\} & \text{if } i < j \end{cases}$$



$$C[i][j] = \begin{cases} 2q_i + p_i + 2q_{i+1} & \text{if } i = j \\ P[i][j] + \min \left\{ \begin{array}{l} q_i + C[i+1][j] \\ C[i][r-1] + C[r+1][j] \text{ for } r \in (i, j) \\ C[i][j-1] + q_{j+1} \end{array} \right\} & \text{if } i < j \end{cases}$$

$$P[i][j] = \begin{cases} q_i + p_i + q_{i+1} & \text{if } i = j \\ \left\{ \begin{array}{l} q_i + p_i + P[i+1][j] \\ \text{or } P[i][r-1] + p_r + P[r+1][j] \text{ for } r \in (i, j) \\ \text{or } P[i][j-1] + p_j + q_{j+1} \end{array} \right\} & \text{if } i < j \end{cases}$$

	0	1	2	3	4	5	6	7	
$k_i$	a	do	i	in	like	not	then	you	
$p_i$	.073	.045	.104	.05	.055	.103	.076	.042	
$q_i$	.001	.113	.107	.006	.001	.02	.081	.122	.001



## Coming up:

*Due **Wed, Nov 16** (end of day)*

*Read Section 6.5*

*(No quiz on Section 6.5)*

*Due **Fri, Nov 18** (end of day)*

*Read Sections 7.(1 & 2)*

*Take quiz*

*Due **Mon, Nov 21** (end of day)*

*Do Project 7.1 (as practice problem)*

*Due **Mon, Nov 28** (end of day) (recommended to be done before break)*

*Read Section 7.3*

*Do Exercises 7.(4,5,7,8)*

*Take quiz*

*Do **Optimal BST** project (suggested by Monday, Nov 21)*