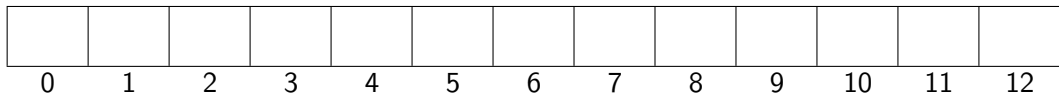


Chapter 6, Hash tables:

- ▶ General introduction; separate chaining (last week Friday)
- ▶ Open addressing (**Today**)
- ▶ Hash table performance (Monday after Thanksgiving)

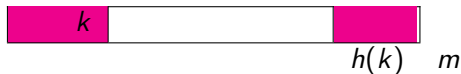
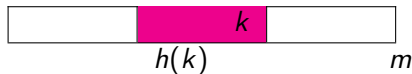
Today:

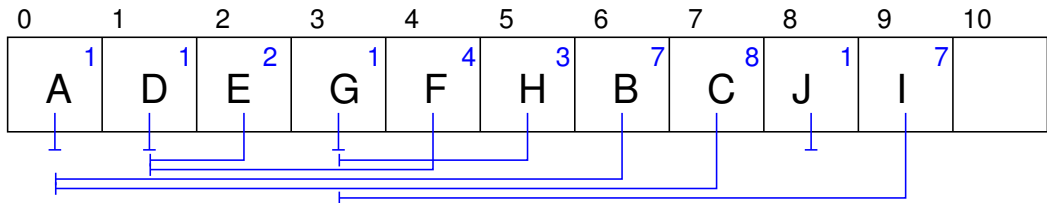
- ▶ Review/finish hash table concepts
- ▶ Basic idea and example of open addressing
- ▶ Terminology, code, and invariant
- ▶ Probing strategies
- ▶ Deletion

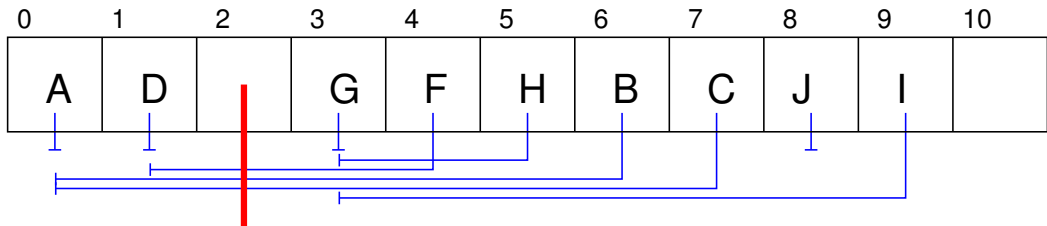


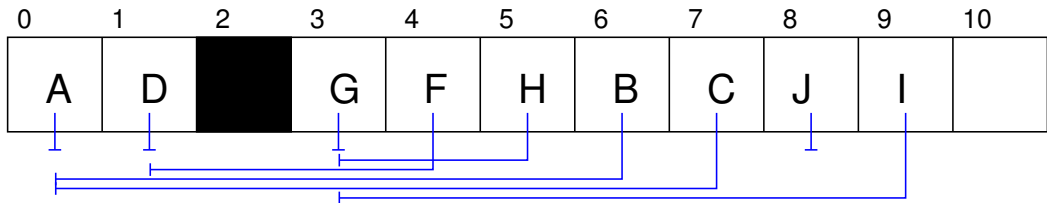
Invariant (Class OpenAddressingHashMap)

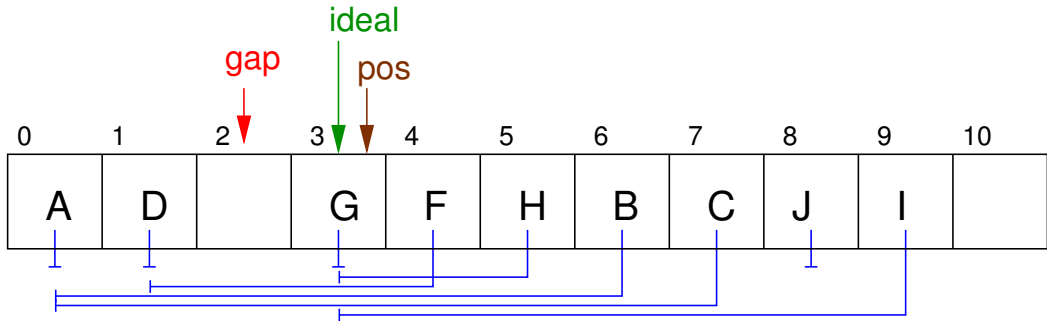
1. The table is not full; there exists $i \in [0, m)$ such that $\text{table}[i] = \text{null}$.
2. There are no breaks in the chain for any key in the table; for all $i \in [0, m)$ such that $\text{table}[i]$ contains key k ,
 - ▶ if $h(k) \leq i$, then for all $j \in [h(k), i]$, $\text{table}[j] \neq \text{null}$;
 - ▶ if $i < h(k)$, then for all $j \in [0, i] \cup [h(k), m)$, $\text{table}[j] \neq \text{null}$.

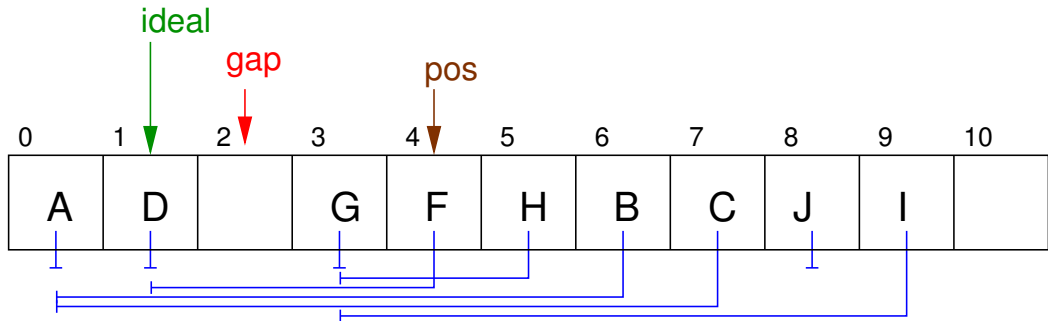


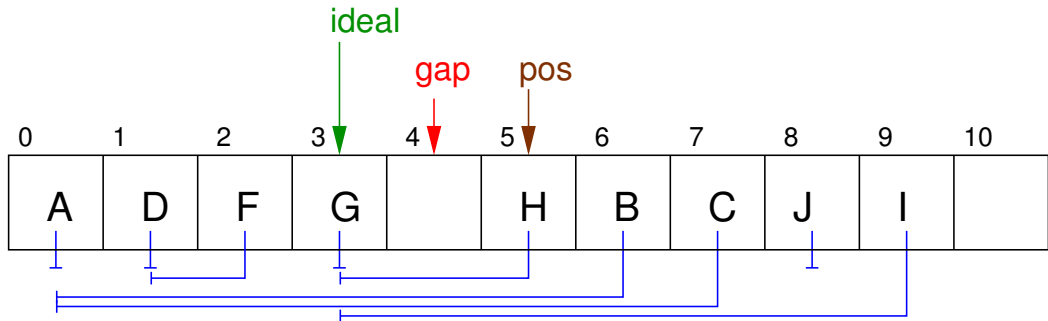


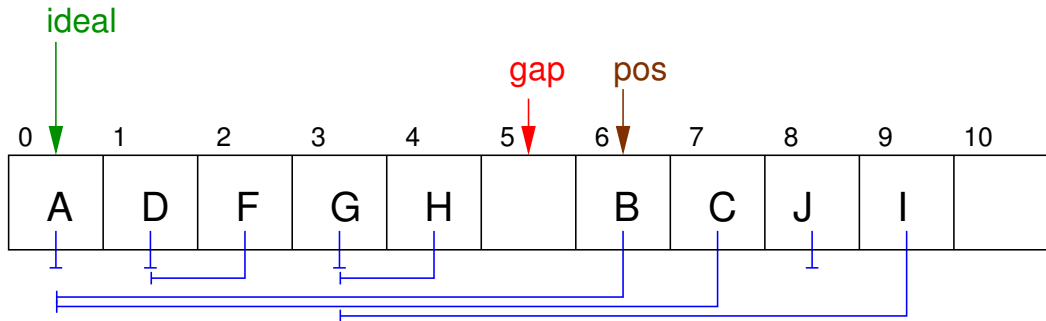


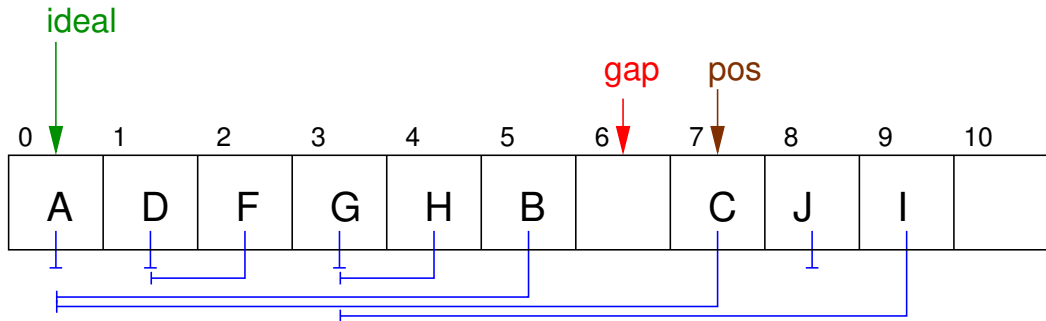


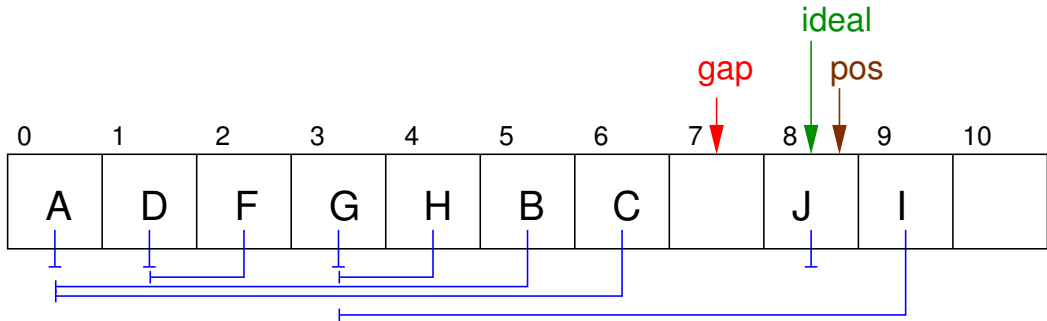


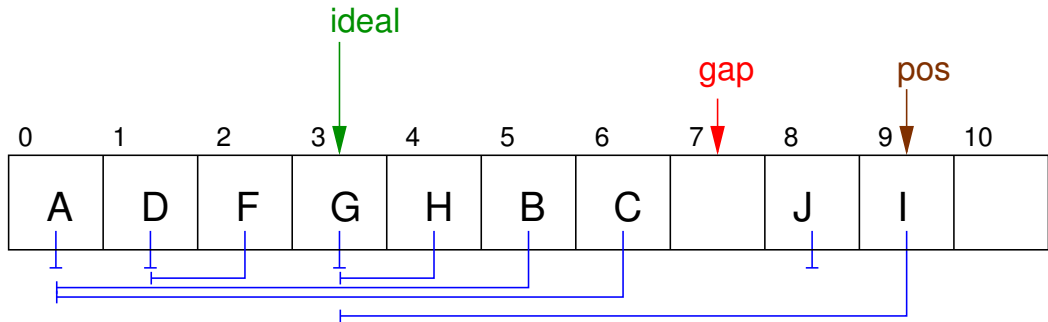


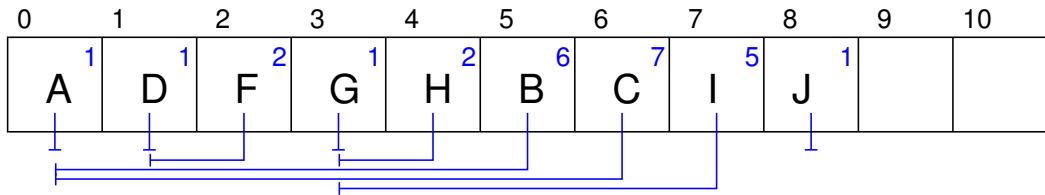


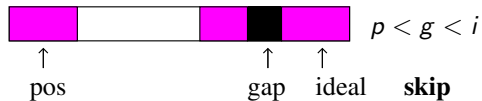
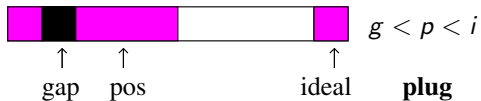
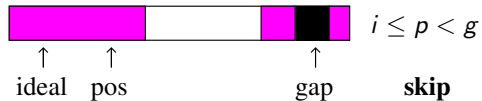
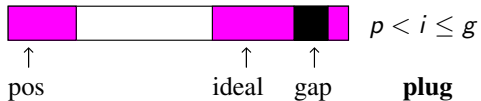
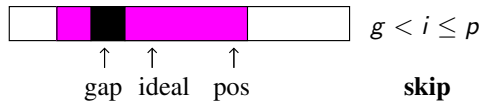
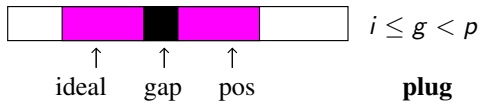












Invariant (Loop of optimized remove in linear probing.)

For all positions $k \in (i, j)$, gap is the only position, if any, between its ideal place ($h(\text{keys}[k])$) and its actual place (k).

Coming up:

Do **Optimal BST** project (*suggested by Today, Nov 21*)

Due **today, Mon, Nov 21** (*end of day*)

Do Project 7.1 (*as practice problem*)

Due **Mon, Nov 28** (*end of day*) (*recommended to be done before break*)

Read Section 7.3

Do Exercises 7.(4,5,7,8)

Take quiz