

## Chapter 5, Binary search trees:

- ▶ Binary search trees; the balanced BST problem (fall-break eve; finished last week Friday)
- ▶ AVL trees (last week Friday and this past Monday)
- ▶ Traditional red-black trees (**Today**)
- ▶ Left-leaning red-black trees (Friday)
- ▶ “Wrap-up” BST (next week Monday)

## Today:

- ▶ Solutions to recent exercises
- ▶ Red-black trees in context
- ▶ Definition and examples
- ▶ Codebase details
- ▶ Cases for put-fixup
- ▶ Analysis



<https://www.beyourownbirder.com>

**Test 1 problem 7. Draw** a weighted, directed, simple graph with five vertices that might require four rounds of relaxations to converge on an SSSP tree starting from 0. **Label** the vertices  $[0, 5)$  and indicate positive weights on the edges. Then **list** the edges in your graph (for example, “(1,3), (4,2), ...”) in an order of relaxations that would require four rounds.

**5.2.** Prove that for any BST, a key that is not already in the tree can be inserted as a leaf.

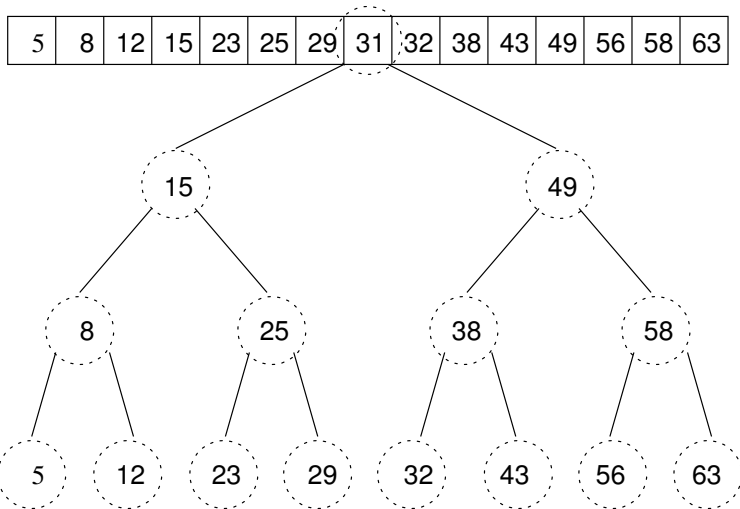
**5.6.** Prove that in a BST, any node with two children has a successor with no more than one child.

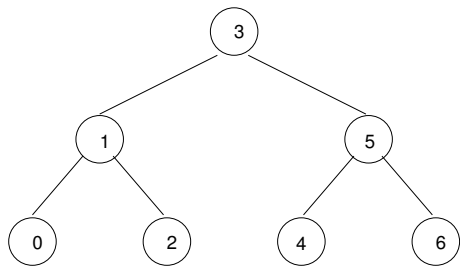
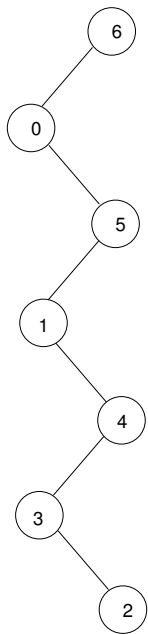
**5.8.** Prove, using structural induction, that for any (*non-empty*) AVL tree there exists a leaf that can be removed without incurring an AVL violation (*and without reducing the height by more than one*).

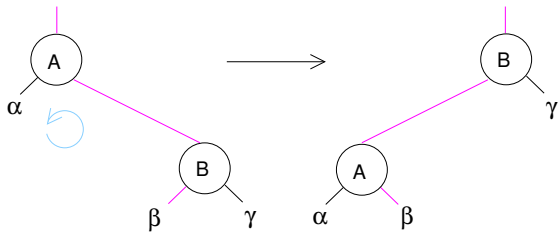
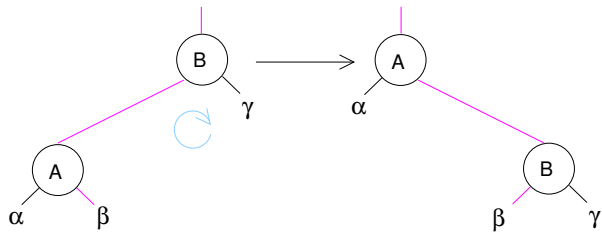
# TRADITIONAL Red-Black Trees

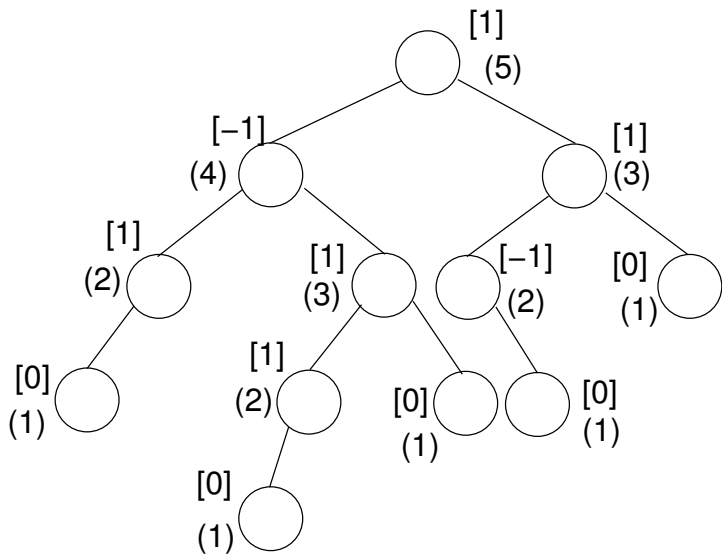


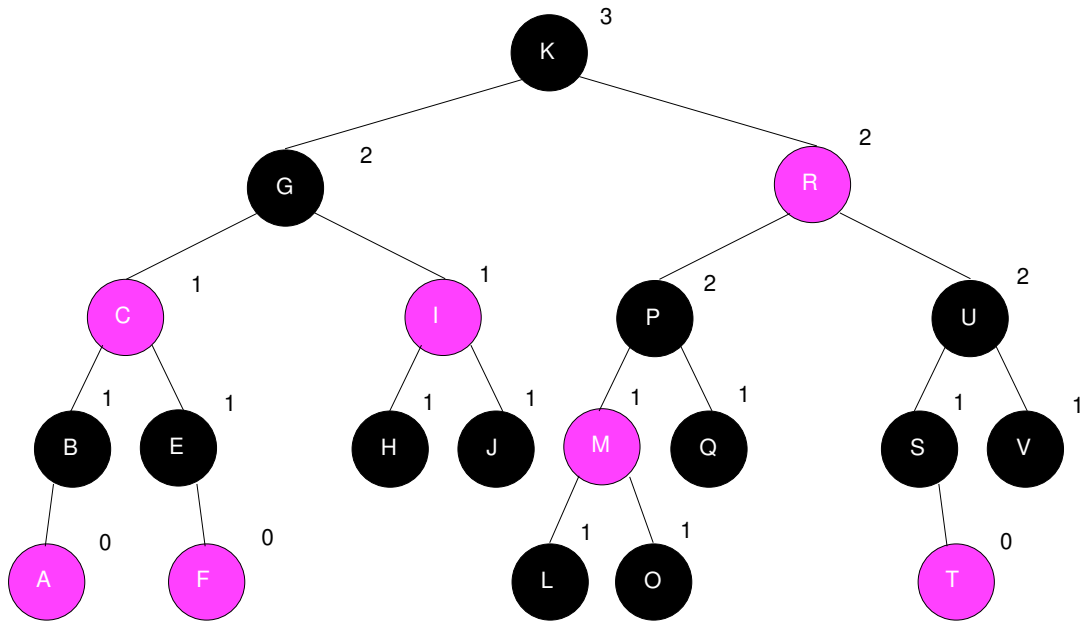
Image swiped from <http://www.cheeseandchickpeas.com/tomato-halloumi-salad/>



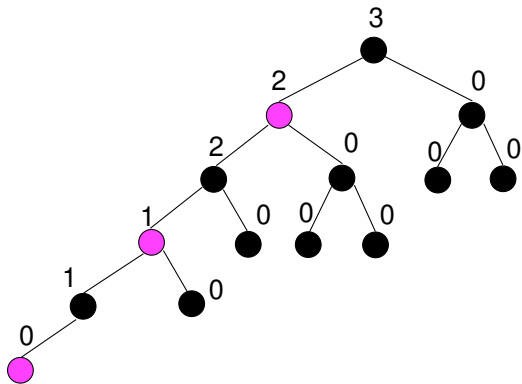
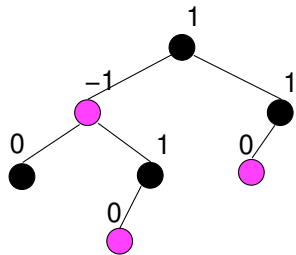






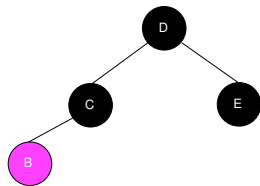
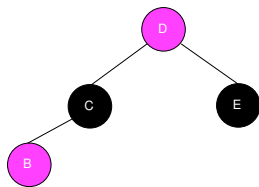
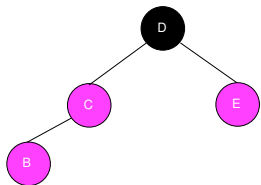
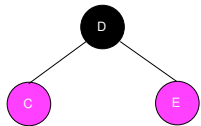


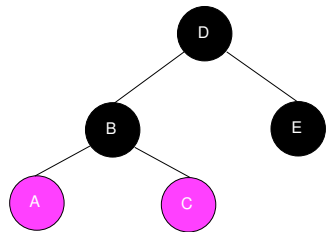
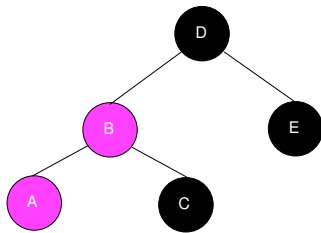
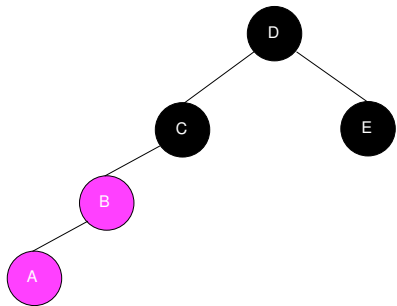




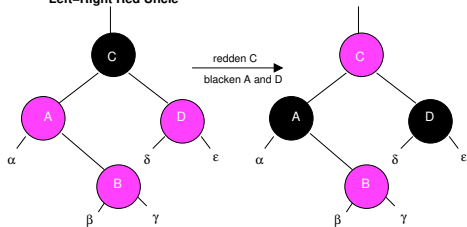
A red-black tree is a binary tree (usually a BST) that is either empty or it is rooted at node  $T$  such that

- ▶  $T$  is either red or black.
- ▶ Both of  $T$ 's children are roots of red-black trees.
- ▶ If  $T$  is red, then both its children are black.
- ▶ The red-black trees rooted at its children have equal blackheight; moreover, the blackheight of the tree rooted at  $T$  is one more than the blackheight of its children if  $T$  is black or equal to that of its children if  $T$  is red.

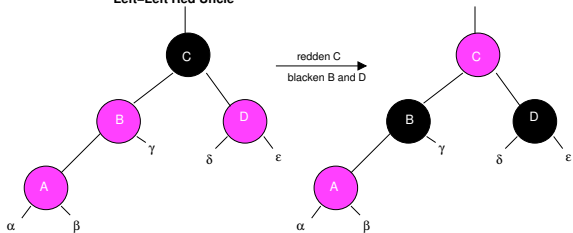




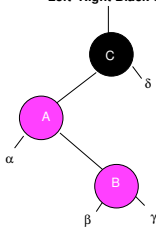
Left-Right Red Uncle



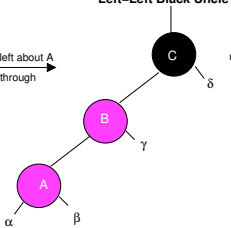
Left-Left Red Uncle



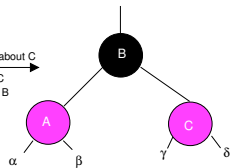
Left-Right Black Uncle



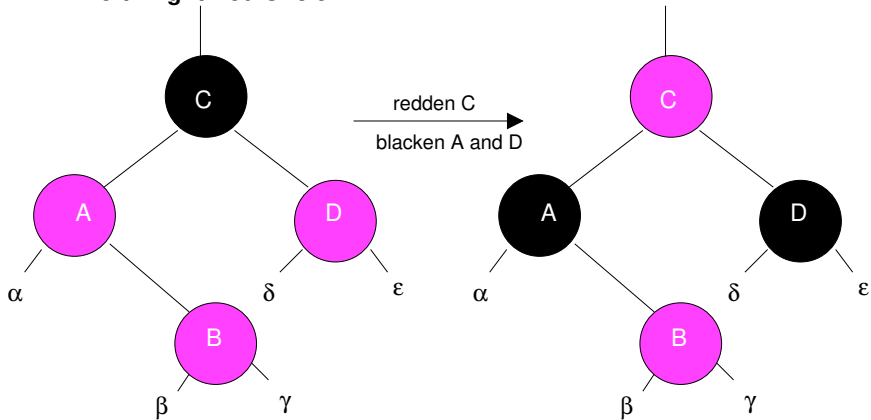
rotate left about A  
fall through



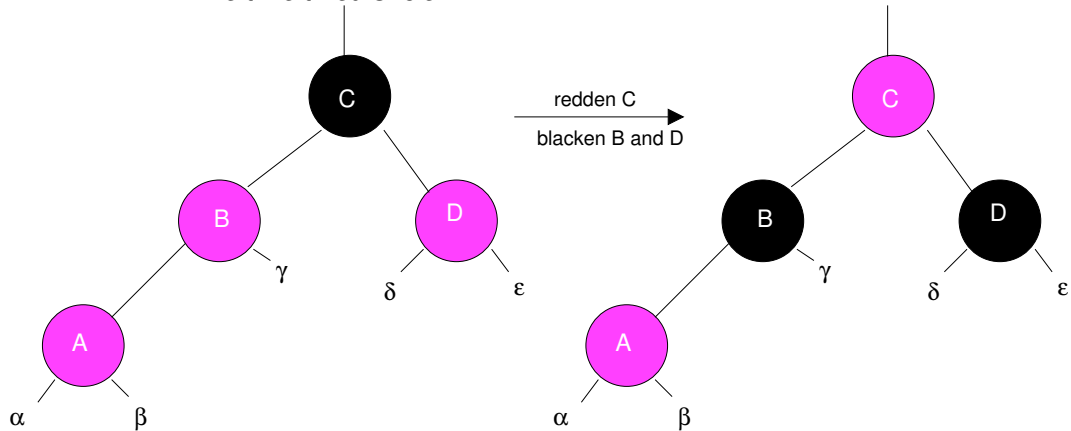
rotate right about C  
redden C  
blacken B



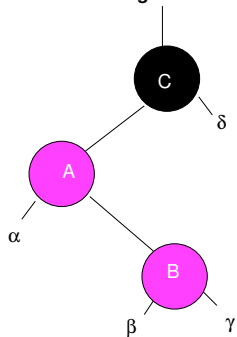
### Left-Right Red Uncle



### Left-Left Red Uncle

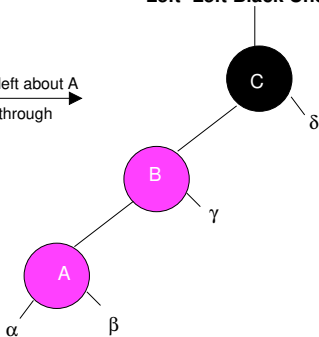


**Left-Right Black Uncle**

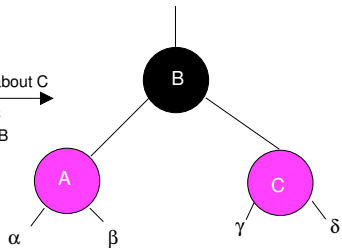


rotate left about A  
fall through

**Left-Left Black Uncle**



rotate right about C  
redden C  
blacken B





**Invariant 26 (Postconditions of RealNode.put() with TradRBBalancer.)** Let  $x$  be the root of a subtree on which `put()` is called and let  $y$  be the node returned, that is, the root of the resulting subtree.

- (a) The subtree rooted at  $y$  has a consistent black height.
- (b) The black height of subtree rooted at  $y$  is equal to the original black height of the subtree rooted at  $x$ .
- (c) The subtree rooted at  $y$  has no double-red violations except, possibly, both  $y$  and one of its children is red.

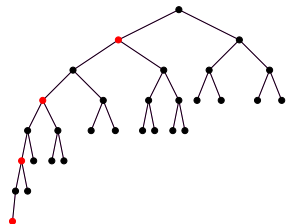
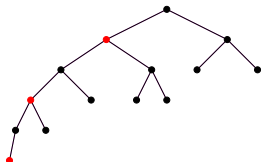
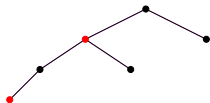
Blackheight

1

2

3

4



Height

2

4

6

8

Nodes

2

6

14

30

## AVL trees

$$h \leq 1.44 \lg n$$

The difference between the longest routes to leaves in the two subtrees is no greater than 1.

Stronger constraint, more aggressive rebalancing, more balanced tree, more work spent rebalancing.

## (Traditional) red-black trees

$$h \leq 2 \lg(n + 2) - 2$$

The longest route to any leaf is no greater than twice the shortest route to any leaf.

Looser constraint, less aggressive rebalancing, less balanced tree, less work spent rebalancing.

## Coming up:

Do **BST rotations** project (*suggested by this past Monday, Oct 24*)

Do **AVL** project (*suggested by Friday, Oct 26*)

Do **Traditional RB** project (*suggested by Friday, Nov 4*)

Due **Mon, Oct 31** (*end of day*) (*but spread it out*)

Read Sections 5.(4-6) [*some parts carefully, some parts skim, some parts optional—see Schoology*]

Do Exercise 5.14

Take quiz