

	Linked list	Sorted array	Heap	Optimized heap
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<code>insert()</code>	$\Theta(1)$	$\Theta(n)$	$\Theta(\lg n)$	$\Theta(\lg n)$
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<code>min()</code>	$\Theta(n)$	$\Theta(1)$	$\Theta(1)$	$\Theta(1)$
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<code>extractMin()</code>	$\Theta(n)$	$\Theta(1)$	$\Theta(\lg n)$	$\Theta(\lg n)$
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<code>decreaseKey()</code>	$\Theta(1)$	$\Theta(n)$	$\Theta(n)$	$\Theta(\lg n)$
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<code>delete()</code>	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$	$\Theta(\lg n)$
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	Linked list	Sorted array	Heap	Optimized heap	Fibonacci heap (amortized)
<code>insert()</code>	$\Theta(1)$	$\Theta(n)$	$\Theta(\lg n)$	$\Theta(\lg n)$	$\Theta(1)$
<code>min()</code>	$\Theta(n)$	$\Theta(1)$	$\Theta(1)$	$\Theta(1)$	$\Theta(1)$
<code>extractMin()</code>	$\Theta(n)$	$\Theta(1)$	$\Theta(\lg n)$	$\Theta(\lg n)$	$\Theta(\lg n)$
<code>decreaseKey()</code>	$\Theta(1)$	$\Theta(n)$	$\Theta(n)$	$\Theta(\lg n)$	$\Theta(1)$
<code>delete()</code>	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$	$\Theta(\lg n)$	$\Theta(\lg n)$
<code>union()</code>	$\Theta(1)$	$\Theta(n)$	$\Theta(n \lg n)$	$\Theta(n \lg n)$	$\Theta(1)$

Fibonacci Heaps and Their Uses in Improved Network Optimization Algorithms

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Abstract. In this paper we develop a new data structure for implementing heaps (priority queues). Our structure, *Fibonacci heaps* (abbreviated *F-heaps*), extends the binomial queues proposed by Vuillemin and studied further by Brown. F-heaps support arbitrary deletion from an n -item heap in $O(\log n)$ amortized time and all other standard heap operations in $O(1)$ amortized time. Using F-heaps we are able to obtain improved running times for several network optimization algorithms. In particular, we obtain the following worst-case bounds, where n is the number of vertices and m the number of edges in the problem graph:

- (1) $O(n \log n + m)$ for the single-source shortest path problem with nonnegative edge lengths, improved from $O(m \log_{(m/n+2)} n)$;
- (2) $O(n^2 \log n + nm)$ for the all-pairs shortest path problem, improved from $O(nm \log_{(m/n+2)} n)$;
- (3) $O(n^2 \log n + nm)$ for the assignment problem (weighted bipartite matching), improved from $O(nm \log_{(m/n+2)} n)$;
- (4) $O(m\beta(m, n))$ for the minimum spanning tree problem, improved from $O(m \log \log_{(m/n+2)} n)$, where $\beta(m, n) = \min \{i \mid \log^{(i)} n \leq m/n\}$. Note that $\beta(m, n) \leq \log^* n$ if $m \geq n$.

Applications of Efficient Mergeable Heaps for Optimization Problems on Trees

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Summary. It is shown how to use efficient mergeable heaps to improve the running time of two algorithms that solve optimization problems on trees.

0. Introduction

A *tree* is a directed graph with a distinguished vertex – the *root* such that there is exactly one directed path from the root to any other vertex. If there is an edge