

$$C[24][2] = \min \begin{cases} \text{use 13-ducat coin, } 1 + C[11][2] \\ \text{don't use 13-ducat coin, } C[24][1] \end{cases}$$

Generalized:

$$C[i][j] = \min \begin{cases} 1 + C[i - D[j]][j] & \text{take the } j\text{th denomination,} \\ & \text{which value } D[j] \\ C[i][j - 1] & \text{skip the } j\text{th denomination} \end{cases}$$

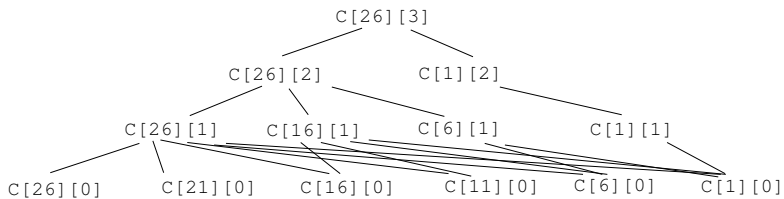
Recursive case (how many of the  $j$ th denomination coins should we take when making change for amount  $i$ ?)

$$C[i][j] = \min_{0 \leq k < x} \{k + C[i - k \cdot D[j]][j - 1]\}$$

$x$  is the exclusive upper bound on how many  $j$ th denomination coins we *could* take, the smallest  $x$  such that  $x \cdot D[j] > i$ .

Base case:

$$C[i][0] = i$$



Recursively characterizing the problem/solution:

$$C[i][j] = \begin{cases} i & \text{if } j = 0 \\ \min_{0 \leq k < x} \{k + C[i - k \cdot D[j]][j - 1]\} & \text{otherwise} \end{cases}$$

Realizing this in pseudocode:

```
For  $i \in [1, amount]$  // loop over all sub-amounts
  For  $j \in [0, D.length)$  // loop over all denomination restrictions
    // find the number of  $j$  coins that minimizes
    // the total number of coins for sub-amount  $i$ 
    // using up to denomination  $j$ 
    Guess the min  $k$  is 0
    For  $k \in [1, x)$ 
      //If  $k$  gets a smaller value than the min so far,
      // make that the new min  $k$ 
    Record the min value as  $C[i][j]$ 
```

$D = [1, 5, 10, 25]$ , *amount* = 26

3	0/0	0/1	0/2	0/3	0/4	0/1	0/2	0/1	0/2	0/3	0/3	1/2				
2	0/0	0/1	0/2	0/3	0/4	0/1	0/2	1/1	1/2	1/3	2/3	2/4				
1	0/0	0/1	0/2	0/3	0/4	1/1	1/2	2/2	2/3	3/4	4/5	5/6				
0	0/0	1/1	2/2	3/3	4/4	5/5	6/6	10/10	11/11	16/16	21/21	26/26				
	0	1	2	3	4	5	6	...	10	11	...	16	...	21	...	26
									<i>i</i>							

$D = [1, 12, 13]$ , *amount* = 24

	2	0/0	0/1	...	0/11	0/1	1/1	...	0/2
<i>j</i>	1	0/0	0/1	...	0/11	1/1	1/2	...	2/2
	0	0/0	1/1	...	11/11	12/12	13/13	...	24/24
		0	1	...	11	12	13	...	24
						<i>i</i>			