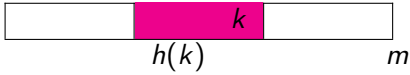
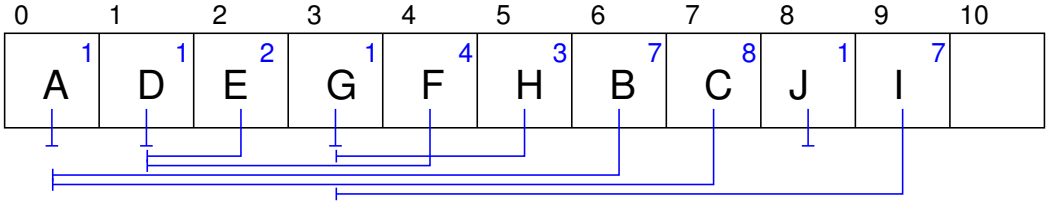
Invariant (Class `OpenAddressingHashMap`)

1. *The table it not full; there exists* $i \in [0, m)$ *such that* $\texttt{table}[i] = \texttt{null}$.
2. *There are no breaks in the chain for any key in the table; for all* $i \in [0, m)$ *such that* $\texttt{table}[i]$ *contains key* k,
   - *if* $h(k) \leq i$, *then for all* $j \in [h(k), i], \texttt{table}[j] \neq \texttt{null}$;
   - *if* $i < h(k)$, *then for all* $j \in [0, i] \cup [h(k), m)$.

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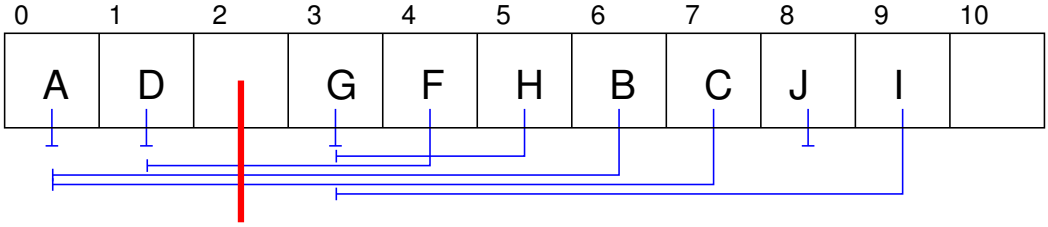
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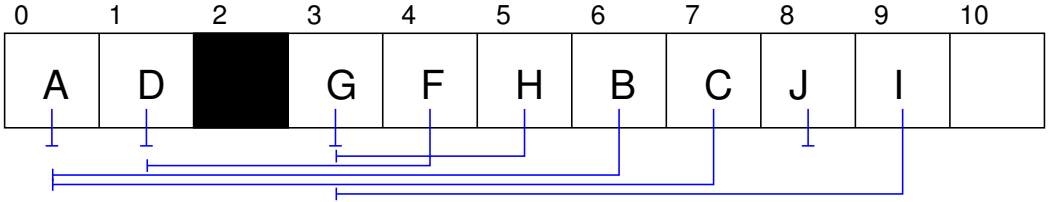
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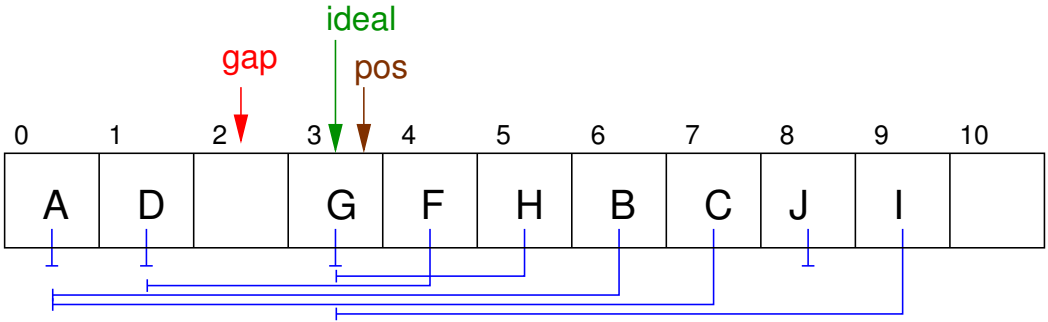
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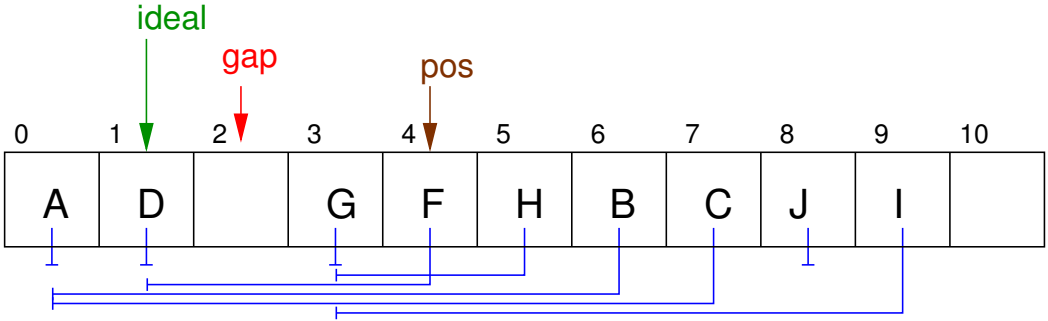
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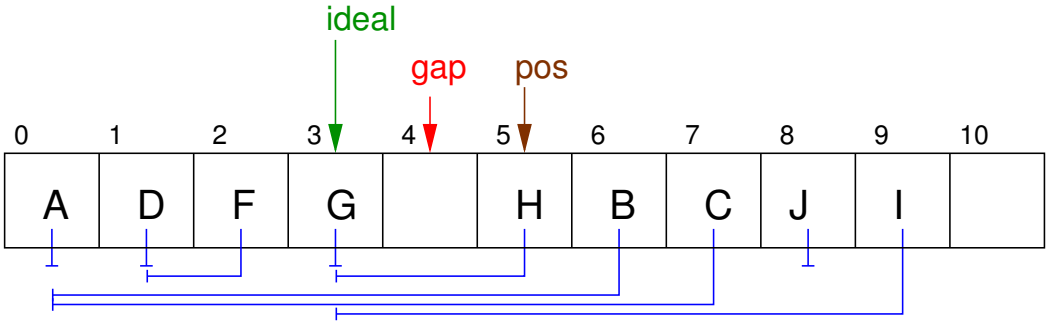
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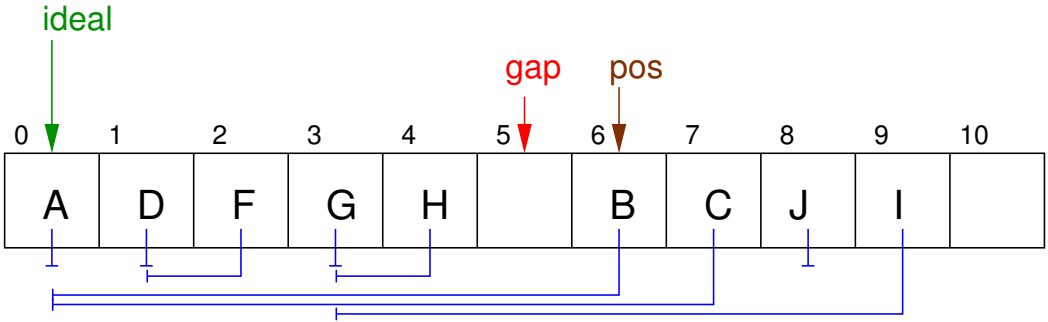
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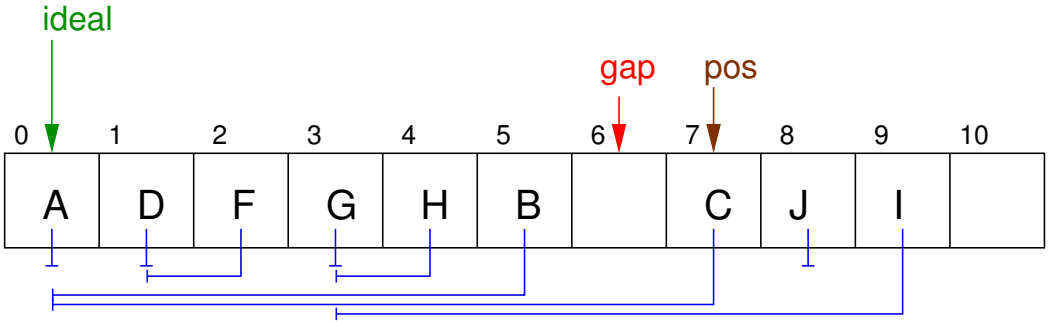
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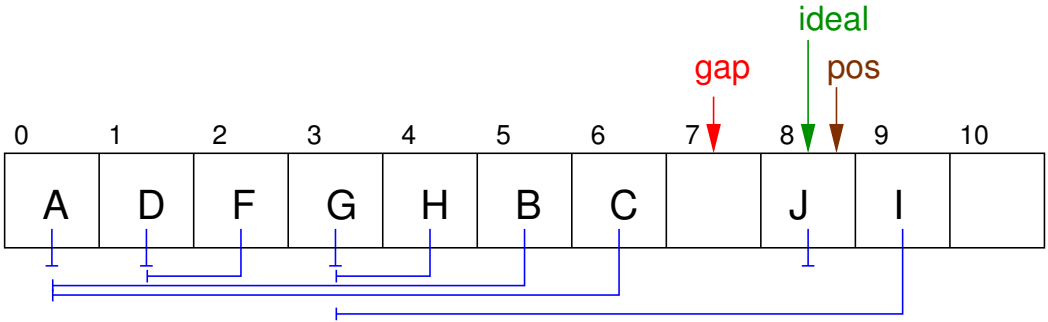
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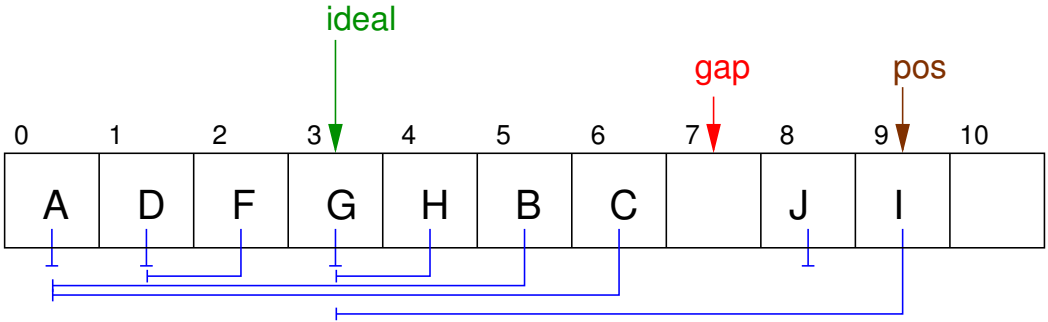
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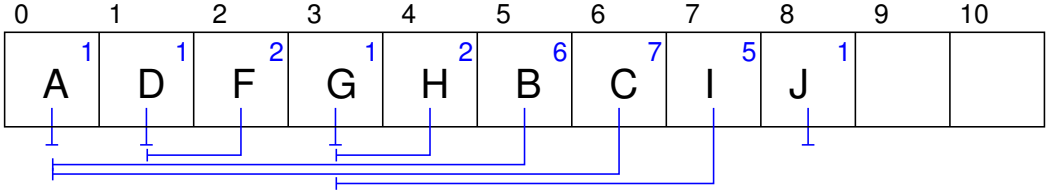
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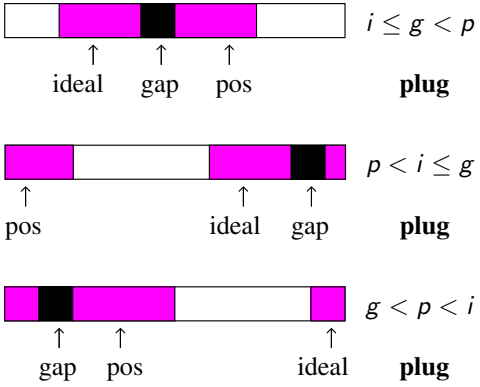

Invariant (Loop of optimized remove in linear probing.)

For all positions $k \in (i, j)$, gap is the only position, if any, between its ideal place ($h(\texttt{keys}[k])$) and its actual place (k).