

Prolegomena unit outline:

- ▶ Algorithms and correctness (Monday and **today**)
- ▶ Algorithms and efficiency (**today** and next week Wednesday and Friday)
- ▶ Abstract data types (the following Wednesday)
- ▶ Data Structures (the following Wednesday and Friday)

Today:

- ▶ Finish check-sorting problem
- ▶ “Binary search” problem
- ▶ Class invariants (`LinkedList`)
- ▶ Start efficiency

What good are invariants?

- ▶ They are a tool for reasoning about the state and progress of an algorithmic process
- ▶ They are a way to explain the meaning of a variable and capture how the variables relate to each other.
- ▶ They help with testing and debugging.
- ▶ They are a means for proving that an algorithm is correct.

Invariant (Class LinkedList)

- (a) $\text{head} = \text{null}$ iff $\text{tail} = \text{null}$ iff $\text{size} = 0$.
- (b) If $\text{tail} \neq \text{null}$ then $\text{tail.next} = \text{null}$.
- (c) If $\text{head} \neq \text{null}$ then *tail* is reached by following $\text{size} - 1$ next links from *head*.

```

def bounded_linear_search(sequence, P):
    a0 found = False
    i = 0
    while not found and i < len(sequence): a1(n + 1)
        a2n found = P(sequence[i])
        i += 1
    if found: a3
        a4 return i - 1
    else :
        a5 return -1

```

$$\begin{aligned}
 T_{bls}(n) &= a_1 + a_2(n + 1) + a_3n + a_4 + \max(a_5, a_6) \\
 &= b_0 + b_1n
 \end{aligned}$$

```

def binary_search(sequence, T0, item):
    c0 low = 0
    high = len(sequence)
    while high - low > 1: c1(lg n + 1)
        c2 lg n mid = (low + high) / 2
        compar = T0(item, sequence[mid])
        if compar < 0: # item comes before mid
            high = mid
        elif compar > 0: # item comes after mid
            low = mid + 1
        else: # item is at mid
            assert compar == 0
            low = mid
            high = mid + 1
    if low < high and T0(item, sequence[low]) == 0: c3
        c4 return low
    else:
        c5 return -1

```

$$\begin{aligned}
 T_{bs}(n) &= c_0 + c_1(\lg n + 1) + c_2 \lg n + c_3 + \max(c_4, c_5) \\
 &= d_0 + d_1 \lg n
 \end{aligned}$$

```

def selection_sort(sequence, T0):
    for i in range(len(sequence)):  $e_0 + e_1 n$ 
        min_pos = i
        min = sequence[i]
        for j in range(i + 1, len(sequence)):  $e_3 n + e_4 \sum_{i=0}^{n-1} (n - i - 1)$ 
            if T0(sequence[j], min) < 0 :  $e_5 \sum_{i=0}^{n-1} (n - i - 1)$ 
                min = sequence[j]
                min_pos = j
            sequence[min_pos] = sequence[i]
            sequence[i] = min

```

$e_2 n$

$$T_{sel}(n) = f_1 + f_2 n + f_3 n^2$$

Coming up:

By class time next week Wednesday:

Do Ex 1.(6 & 7)

Take quiz

By midnight next week Wednesday:

Read Section 1.2

For next week Friday:

Read Sections 1.(3 & 4)

Do practice problems 1.(27 & 28) and 1.(42 & 43)

Take quiz