

## Chapter 6, Hash tables:

- ▶ General introduction; separate chaining (last Friday)
- ▶ Open addressing (**Today**)
- ▶ Hash functions (Wednesday)
- ▶ Perfect hashing (Monday after next)
- ▶ Hash table performance (Wednesday after next)

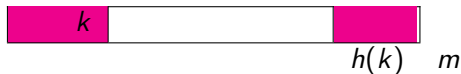
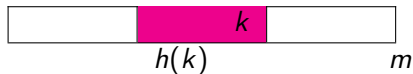
## Today:

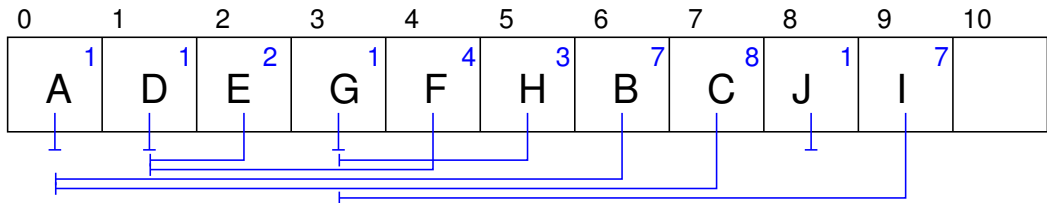
- ▶ Review/finish hash table concepts
- ▶ Basic idea and example of open addressing
- ▶ Terminology, code, and invariant
- ▶ Probing strategies
- ▶ Deletion

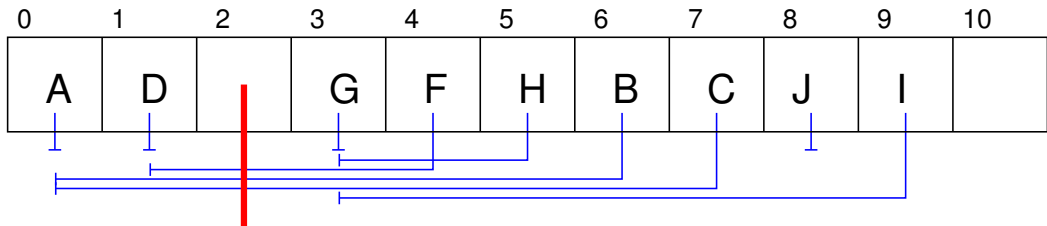


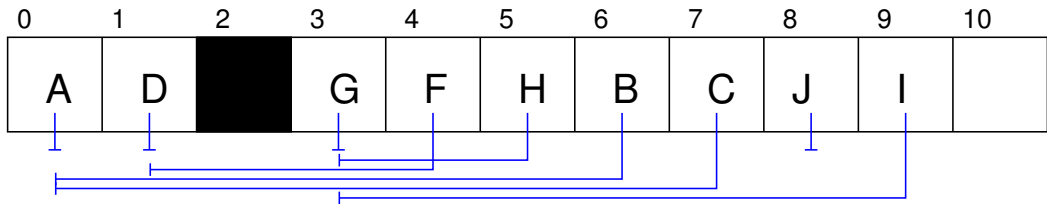
## Invariant (Class OpenAddressingHashMap)

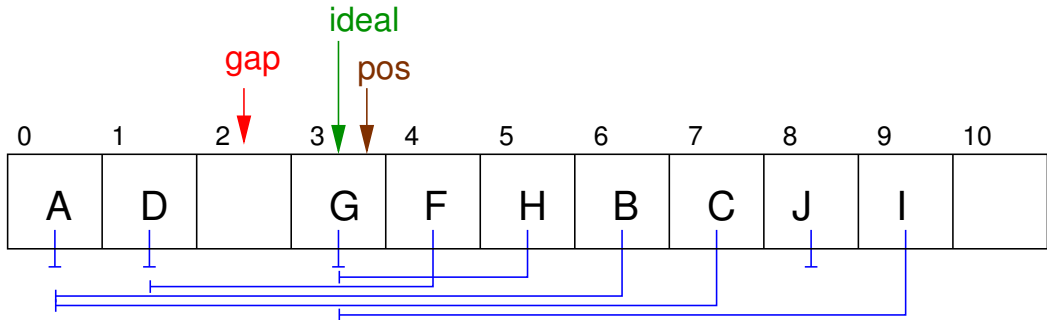
1. The table is not full; there exists  $i \in [0, m)$  such that  $\text{table}[i] = \text{null}$ .
2. There are no breaks in the chain for any key in the table; for all  $i \in [0, m)$  such that  $\text{table}[i]$  contains key  $k$ ,
  - ▶ if  $h(k) \leq i$ , then for all  $j \in [h(k), i]$ ,  $\text{table}[j] \neq \text{null}$ ;
  - ▶ if  $i < h(k)$ , then for all  $j \in [0, i] \cup [h(k), m)$ ,  $\text{table}[j] \neq \text{null}$ .

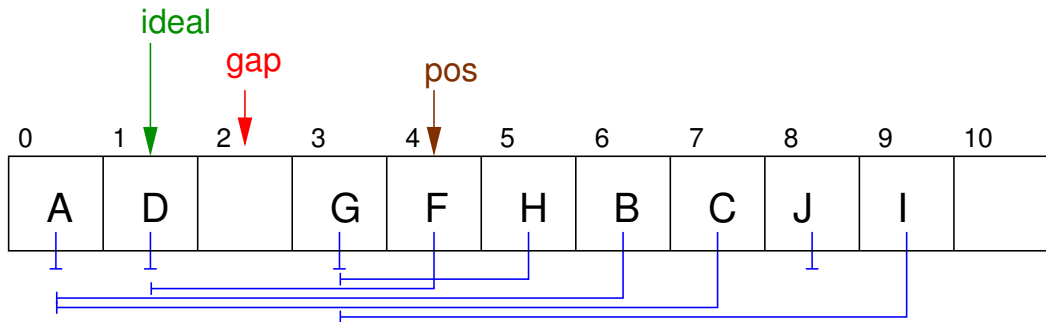




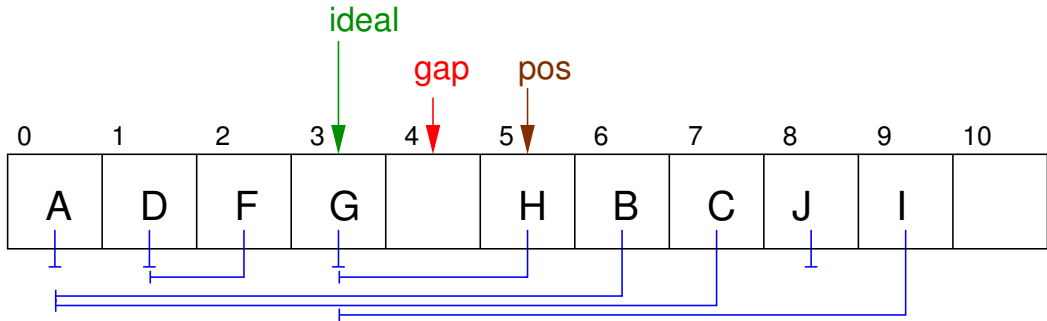


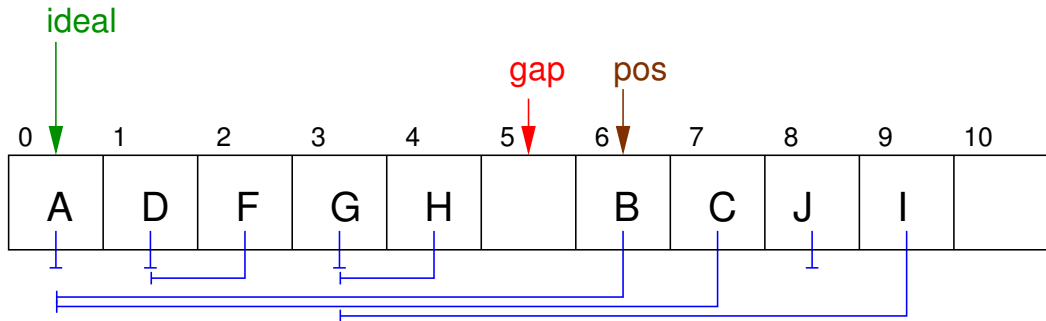


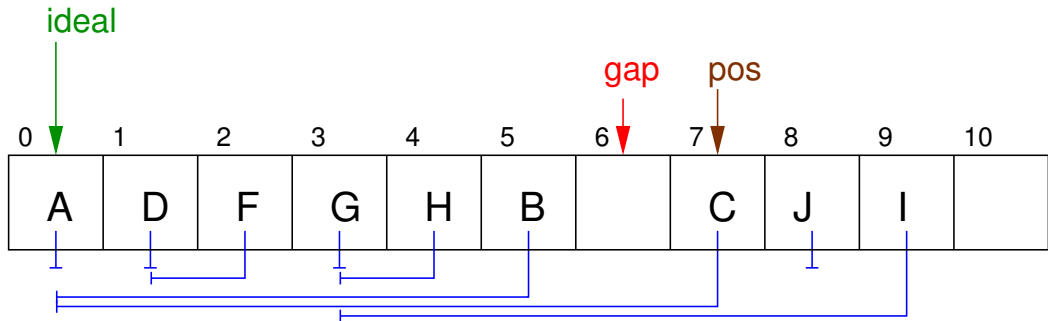


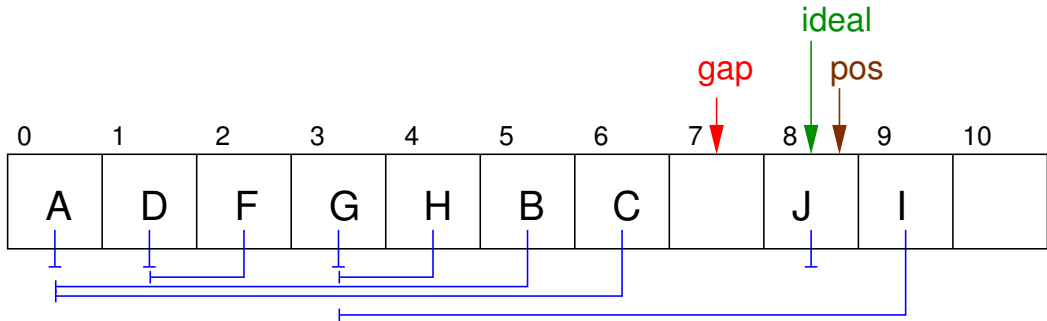


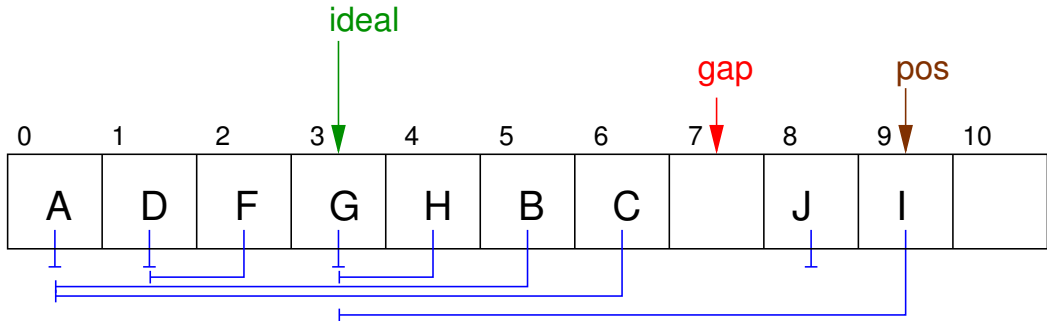


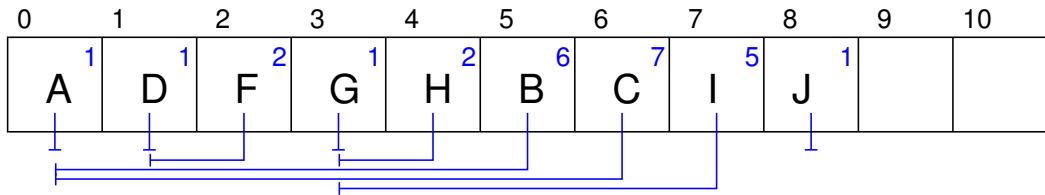


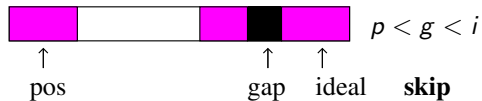
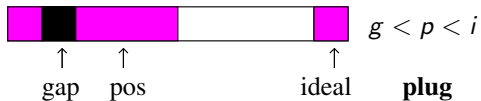
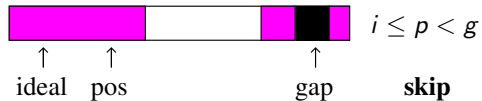
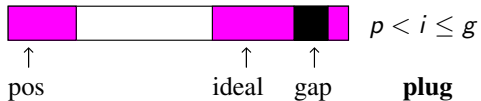
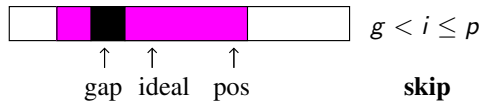
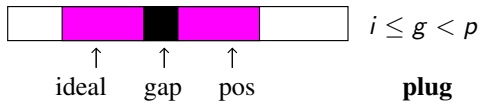












## Invariant (Loop of optimized remove in linear probing.)

For all positions  $k \in (i, j)$ , gap is the only position, if any, between its ideal place ( $h(\text{keys}[k])$ ) and its actual place ( $k$ ).

## Coming up:

Do **Optimal BST** project (suggested by this past Friday, April 8)

Do **Open addressing with linear probing** project (suggested by Monday, Apr 18)

**Due Tues, Apr 12**

*Do practice problem, recreating separate chaining example*

*Read Section 7.3*

*Take quiz*

**Due Mon, Apr 18**

*Read Sections 7.(4 & 5)*

*(No practice problems or quiz)*