#### Chapter 4, Graphs:

- Concepts and implementation (week-before Friday)
- Traversal (last week Monday)
- Minimum spanning trees (last week Wednesday and Friday)
- ► Single-source shortest paths (this week Wednesday and Friday)
- Review for test (next week Monday)

#### "Today" (Wednesday and Friday):

- (MST loose ends)
- ► The SSSP problem
- General concepts for SSSP algorithms
- ▶ The most unlucky graph for SSSP
- ► The Bellman-Ford algorithm plus analysis
- Dijkstra's algorithm plus analysis

### Minimum Spanning Tree Problem

Given a weighted, undirected graph, find the tree with least-total weight that connects all the vertices, if one exists.

## Single-Source Shortest Paths Problem

Given a weighted directed graph and a source vertex, find the tree comprising the shortest paths from that source to all other reachable vertices.

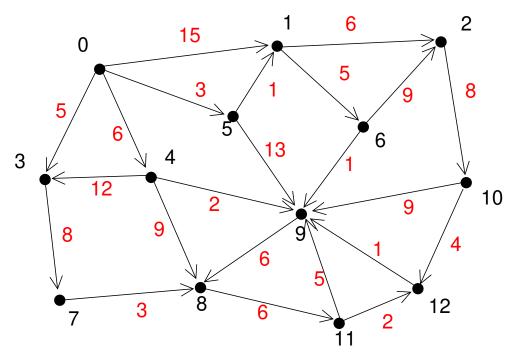
- Both are defined for weighted graphs
- Both produce trees as a result
- ▶ Both minmize by weight
- For each we have two algorithms

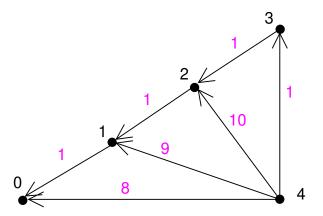
Input is only a graph
Problem usually is described on an undirected graph
Goal is to minimize total weight
There is no clear winner between the algorithms

Input is a graph and a starting point Problem usually is described on a directed graph Goal is to minimize weight on each path

One algorithm is clearly more efficient

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Let X be the set of vertices whose distance bounds are correct, that is,  $v \in X$  if  $\mathtt{distances}[v]$  is the total weight of the shortest path from s to v. For a single-source shortest path algorithm to be correct, all vertices reachable from s are in set X at termination, and if s vertices are reachable, this implies s vertices that have been removed from the priority queue. Our intent is that s vertices have correct distance bound at the time they are removed from the priority queue, though at any point there may also be some correct ones still in the priority queue. We claim

# Invariant (Main loop of Dijkstra's algorithm)

Let X and Y be as defined above.

- (a)  $Y \subseteq X$ .
- (b) If v is the vertex in the priority queue with least distance bound, then  $v \in X$ .
- (c) |Y| is the number of iterations completed.

#### Coming up:

Do **MST** project (due Wednesday, Feb 21) Do **SSSP** project (due Friday, Mar 1)

Due Fri, Feb 23 (end of day)
Read Section 4.5
Do Exercises 4.(50, 51, 59)
Take SSSP quiz